

## A NEW PROCESS FOR SOLVING LINEAR FRACTIONAL PROGRAMMING PROBLEM

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### *Abstract*

*There is a need for generalizing the simplex technique for linear programming to the ratio of linear functions or the ratio of quadratic functions and in such a situation, all the problems are fragments of a general class of optimization problems, termed in the literature as fractional programming problems. The linear fractional programming problem arise when there appears a necessity to optimize the efficiency in other activities also, for example, profit gained by company per unit of expenditure of labor, cost of production per unit of produced goods etc. Nowadays, because of deficit of natural resources, the use of such specific criteria becomes more and more topical and relevant. Linear Fractional problems arises in management decision making, research, finance, production, Health care and Hospital Planning and transportation etc., In this paper, a new process has been introduced for solving linear fractional programming problem. The main purpose of this method is to reduce the number of iterations in order to save time. By using the numerical examples, we suggest that the optimality of this method arises in a few numbers of iteration than the existing Simplex method for solving linear fractional programming problem. Also, in certain cases, the same number of iterations is maintained for the optimality.*

**Keywords:** Basic feasible solution, optimum solution, alternative simplex method, linear programming problem, linear fractional programming problem.

### INTRODUCTION

Linear-fractional programming (LFP) is a generalization of [linear programming](#) (LP) in which the objective function in a linear program is a [linear function](#) whereas the objective function in a linear-fractional program is a ratio of two linear functions. This field of LFP was developed by Hungarian mathematician B. Matros in 1960. Linear Fractional Programming Problems is studied by many authors charnes et al.(1962) reformulate Linear Fractional Programming Problem into a Linear Programming Problem. Among them M. Karade and P.G.Kumbhare gives the alternative approach for the simplex method (1980).Linear fraction programs and Singh C.(1981) discussed a useful study about the optimality condition in Fractional Programming. An integer solution Fractional Programming Problem was discussed by S.C.Sharma ( 2011 ) . Furthermore, Many researchers discussed about the alternative solutions of simplex method. Optimum solution of simplex method and Conventional Simplex Method, An Alternative Approach was given by N.W.Khobragade.(2013). Sophia Porchelvi R and Anna Sheela A introduced a new method for Linear Multi – Objective Fractional Transportation Problem. This paper is organized as follows. In section 2, Mathematical formulation for linear fractional programming problem is given along with some basic definitions. An algorithm for solving linear fractional programming problem using conventional simple method is developed in section 3 and it is verified by the numerical example. The last section draws some concluding remarks.

**MATHEMATICAL FORMULATION**

**1** The general linear programming problem in the form

$$\text{Maximize } Z = c_1 x_1 + c_2 x_2 + \dots + c_n x_n$$

Subject to the constraints

$$a_{i1} x_1 + a_{i2} x_2 + \dots + a_{in} x_n = b_i$$

$$x_1, x_2, \dots, x_n \geq 0$$

**2** A linear fractional programming problem:

$$\text{Maximize } Z = \frac{c x + \alpha}{d x + \beta}$$

Subject to the constraints

$$Ax \leq b, x \geq 0.$$

Where i) A is an  $n \times m$  matrix,  $A = [a_{ij}]$  ( $i = 1, \dots, m; j = 1, \dots, n$ )

ii) x, c and d are  $(n \times 1)$  vectors, b is an  $(m \times 1)$  vectors and

iii)  $\alpha, \beta$  are scalars.

**DEFINITIONS**

**(i) Degenerate Solution**

A basic solution to the system is called Degenerate, if one or more of the basic variables vanish.

**(ii) Basic Feasible Solution**

A feasible solution to a Linear Programming Problem, which is also a basic solution to the problem is called a Basic Feasible Solution.

**(iii) Optimal Solution**

A basic feasible solution  $x_B$  to the Linear Programming Problem

$$\text{Maximize } z = c x$$

Subject to the constraints:  $A x \leq b, x \geq 0$ , is said to be an Optimal Solution, if all  $c_j - z_j \leq 0$ .

**NEW ALGORITHM FOR SOLVING LINEAR FRACTIONAL PROGRAMMING PROBLEM**

**STEP 1:**

First, observe whether all the right-side constants of the constraints are non-negative. If not, it can be changed into positive value on multiplying both the sides of the constraints by (-1).

**STEP 2:**

Converts the inequality constraints to equations by introducing the non-negative slack or surplus variables. The coefficients of slack or surplus variables are always taken zero in the objective function.

**STEP 3:**

Let  $x_B$  be the initial basic feasible solution of the given problem such that

$$B x_B = b$$

$$x_B = b B^{-1} \quad \text{where } B = (b_1, b_2, \dots, b_r, b_s, \dots, b_m).$$

$$\text{Further suppose that, } Z_1 = c_B x_B + \alpha \text{ and } Z_2 = d_B x_B + \beta$$

Where  $c_B$  and  $d_B$  are the vectors having their components as the coefficients associated with the basic variables in the numerator and denominator of the objective function respectively.

**STEP 4:**

Compute the evaluations  $c_j - z_j$  ( $j=1,2,\dots,n$ ) by using the relation

$c_j - z_j = c_j - c_B y_j$ , Where  $Y_j = B^{-1}a_j$ .

Also compute  $\frac{c_j - z_j}{\sum y_j}$

Case (i) If  $y_{ij} > 0$ , then Go to Step 6

Case (ii) If  $y_{ij} < 0$ , then Go to Step 6

Case (iii) If  $y_{ij} = 0$ , then that column is undefined and we don't consider that column for further process.

Similar, computation can be made for denominator also.

**STEP 5:**

Now, compute the net evaluation  $\Delta_j$  for each variable  $x_j$  by the formula

$$\Delta_j = z_2 \left( \frac{c_j - z_j}{\sum y_j} \right) - z_1 \left( \frac{d_j - z_j}{\sum y_j} \right)$$

**STEP 6:**

Case (i) If all  $\Delta_j \leq 0$ , then the initial basic feasible solution  $X_B$  is an optimum basic feasible solution.

Case (ii) If at least one  $\Delta_j > 0$ , go to the next step.

**STEP 7:**

If there is more than one positive  $\Delta_j$ , then choose the most positive of them. Let it be  $\Delta_r$  for some  $j = r$ .

Case (i) If all  $y_{ir} \leq 0$ , ( $i = 1, 2, \dots, m$ ), then there is an unbounded solution to the given problem.

Case (ii) If at least one  $y_{ir} > 0$ , ( $i = 1, 2, \dots, m$ ), then the corresponding vector  $y_r$  enters the basis  $y_B$ .

**STEP 8:**

Compute the ratios  $\left\{ \frac{x_{Bi}}{y_{ir}}, y_{ir} > 0, i = 1, 2, \dots, m \right\}$  and choose the minimum of them. Let the minimum of these ratios be  $\frac{x_{Bk}}{y_{kr}}$ . Then the vector  $y_k$  will leave the basis  $y_B$ . The common element  $Y_{kr}$ , which is in the  $k^{\text{th}}$  row and the  $r^{\text{th}}$  column is known as the leading element (or pivotal element) of the table.

**STEP 9:**

Convert the leading element to unity by dividing its row by the leading element itself and all other elements in its column to zeroes by making use of the relations:

$$Y_{ij} = y_{ij} - \frac{y_{kj}}{y_{kr}} y_{ir}, \quad i=1,2,\dots,m, i \neq k$$

$$Y_{kj} = \frac{y_{kj}}{y_{kr}} \quad j=1, 2, \dots, n.$$

**Step 10:**

Go to Step 4 and repeat the computational procedure until either an optimum solution is obtained or there is an unbounded solution.

**NUMERICAL EXAMPLE**

Find the integer solution of the following fractional linear programming problem:-

$$\text{Maximize } z = \frac{2x_1 + x_2}{3x_1 + x_2 + 6}$$

Subject to

$$5x_1 + 3x_2 \leq 6$$

$$7x_1 + x_2 \leq 6$$

$$x_1, x_2 \geq 0$$

**SOLUTION** After adding slack variables  $x_3$  and  $x_4$ , the simplex table gives

**Table 1**

			$c_j$	2	1	0	0	Min. Ratio
			$d_j$	3	1	0	0	$\theta$
$d_B$	$C_B$	$X_B$	B	$x_1$	$x_2$	$x_3$	$x_4$	
0	0	$x_3$	6	5	3	1	0	2
0	0	$x_4$	6	7	1	0	1	6
$\sum y_j$				12	4	1	1	
$c_j - z_j$				2	1	0	0	
$\frac{c_j - z_j}{\sum y_j}$				1/6	1/4	0	0	
$d_j - z_j$				3	1	0	0	
$\frac{d_j - z_j}{\sum y_j}$				1/4	1/4	0	0	
$\Delta_j$				1	3/2	0	0	

Therefore,  $x_1$  enters into the basis and  $x_3$  leaves the basis

**Table 2**

			$c_j$	2	1	0	0
			$d_j$	3	1	0	0
$d_B$	$C_B$	$X_B$	B	$x_1$	$x_2$	$x_3$	$x_4$
1	1	$x_2$	2	5/3	1	1/3	0
0	0	$x_4$	4	16/3	0	-1/3	1
$\Delta_j$				0	0	-	0

Here  $x_1 = 0, x_2 = 2$  and  $Z = 1/4$ . Since all  $\Delta_j \leq 0$ , the current solution is optimal.

**NUMERICAL EXAMPLE**

$$\text{Maximize } z = \frac{3x_1 + x_2}{2x_1 + 7}$$

Subject to

$$x_1 + 2x_2 \leq 3$$

$$3x_1 + 2x_2 \leq 6$$

**SOLUTION**

**TABLE 1**

			$c_j$	6	5	0	0	Min. Ratio
			$d_j$	2	0	0	0	$\theta$
$d_B$	$c_B$	$x_B$	B	$x_1$	$x_2$	$x_3$	$x_4$	
0	0	$x_3$	3	1	2	1	0	3
0	0	$x_4$	6	3	2	0	1	2
$\sum y_j$				4	4	1	1	
$c_j - z_j$				6	5	0	0	
$\frac{c_j - z_j}{\sum y_j}$				3/2	5/4	0	0	
$d_j - z_j$				2	0	0	0	
$\frac{d_j - z_j}{\sum y_j}$				1/3	0	0	0	
$\Delta_j$				21/2	35/4	0	0	

Here,  $x_1$  enters into the basis and  $x_4$  leaves the basis.

**TABLE 2**

			$c_j$	6	5	0	0	Min. Ratio
			$d_j$	2	0	0	0	$\theta$
$d_B$	$c_B$	$x_B$	B	$x_1$	$x_2$	$x_3$	$x_4$	
0	0	$x_3$	1	0	4/3	1	-1/3	3/4
2	6	$x_1$	2	1	2/3	0	1/3	3
$\Delta_j$				0	27/2	0	-	

Therefore,  $x_2$  enters into the basis and  $x_3$  leaves the basis.

**TABLE 3**

			$c_j$	6	5	0	0
			$d_j$	2	0	0	0
$d_B$	$c_B$	$x_B$	B	$x_1$	$x_2$	$x_3$	$x_4$
0	5	$x_2$	3/4	0	1	3/4	-1/4
2	6	$x_1$	3/2	1	0	-1/2	1/2
$\Delta_j$				0	0	-81	-19

Here  $x_1 = 3/2$ ,  $x_2 = 3/4$  and  $Z = Z_1 / Z_2 = 51 / 40$ . Since all  $\Delta_j \leq 0$ , therefore the current

solution is optimal.

### CONCLUSION

In this paper, an alternative simplex method has been introduced for the solving linear fractional programming problem. By the numerical examples it has been verified and concluded that the iterations required for optimum solution of this method are less or almost equal as compared to the existing method for solving the linear fractional programming problem.

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